Analysis and Implementation of Adequate Database Management System in Wireless Sensor Networks

Mohd Muntjir[1], Mohd Rahul[2] and Hussain Abu Sorrah[3]
[1,2,3]College of Computers and Information Technology, Taif University, Taif, Kingdom of Saudi Arabia

***Abstract*—Wireless sensor networks are an emerging area of research interest and a network of distributed sensors grouped together to monitor physical or environmental conditions, like pressure, temperature, sound etc. The development of wireless sensor networks was first motivated by military applications; today such networks are used in several industrial, non-industrial and consumer applications, such as industrial process monitoring and control, machine observation, health monitoring, etc. By framework sensor networks as virtual databases, we can offer a nonprocedural-programming interface suitable to data management system. We squabble here that in order to attain energy efficient and useful completion, query-processing operators should be implementing within the sensor network and that estimated query results would play a key role in network system. We study that in network implementations of database operators need novel data centric routing mechanism, as well as a reassessment of conventional network and database interface layering. Wireless sensor networks are presently getting considerable concentration due to their unrestrained prospective.**

***Keywords*—Wireless Network, Sensor Network, Sensor Node, Query Processing and Data Management in Wireless Sensor Networks**

I. INTRODUCTION

A sensor network is viewed as a distributed database that accumulates physical capacity about the environment indexes them and then serves queries from user. Each sensor node normally generates a stream of data items that are obtained from the sensing devices on the node [6]. Wireless sensor networks have acknowledged important recent concentration in both the networking and operating systems community [23], [25]. Anticipating the development of such devices, recent work has also begun exploring possible applications of sensor networks for monitoring diverse infrastructures. The examples of such applications such as monitoring in edifice energy usage for preparation energy maintenance [1]; military and civilian surveillance [21]; monitoring of natural habitats with a view to understanding ecosystem atmosphere [6]; and data assembling in instrumented learning framework for children [35]; and measuring variations in local salinity levels in riparian atmosphere [36]. Particularly, the energy cost of contact is expected to be considerably higher than the cost of local calculation [30, 23]Sensor networks are best intended in a data centric process: the low-level communication primitives in these networks are considered in terms of named data rather than the node identifiers used in conventional networked communications [51]. In this paper, we investigate challenges in realizing this framework view of the sensor network as a database system. Particularly, this view allows users to concern database queries to one or more nodes within the sensor networks [52]. These queries can be one-shot relational queries with a fixed answer set or ongoing permanent queries that create an unrestrained stream of consequences. In a typical situation, users can recover information of interest from a WSN by injecting queries and gathering results from the so-called base stations which behave as an interface between users and the network. It is also intended that sensor networks will eventually be connected to the Internet through which inclusive information sharing becomes realistic.

Furthermore, we propose that a sensor network database or a sensornet database should be architected on two significant thoughts [13].

The first is in network implementations of primal database query operators such as grouping, aggregation, and joins. By in network we mean group communication and routing protocols which together with probable dispensation at intermediary nodes that execute each operator in an application sovereign method [14].

Second, dissimilar the strict semantics connected with customary data models and query languages, we bicker for comforting the semantics of database queries to allow estimated outcome. This recreation enables energy-efficient implementations even given the predictable high level of network dynamics [15]. Hence, it is not only expedient but certainly more precise to present estimated semantics and represent a spectrum of tradeoffs between brief and precise communication system. As we converse underneath, numerous pieces of prior work on online sampling and estimation in the database neighborhood are appropriate in this perspective.

II. ENVIRONMENT AND STRUCTURE OF SENSOR NODES

The sensor node comprises of sensors, microcontrollers, and RF transceiver. It is often compelled by a battery or energy collecting structure. The sensor produces analogue signals, and an ADC transforms the signals. The microcontroller implements a series of algorithms to develop the data. All data will be amassed in the microcontroller and transmit through a unified RF transceiver.

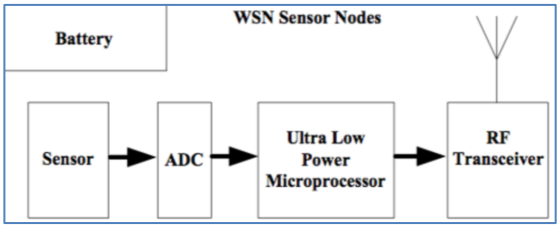

Fig. 1. The system block diagram for Wireless Sensor Network nodes

The data produced in a sensor network is basically the readings of the sensing devices on the nodes, and can be demonstrated as relational data streams [25] is used to managing aggregation in in-network approach. Further, Tiny DB ropes metadata executive, in-network insistent storage.

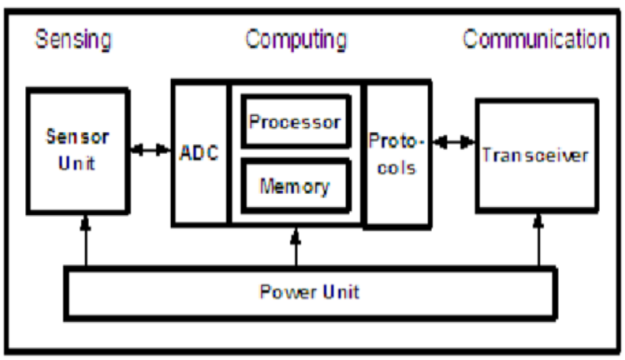

Fig. 2. The typical architecture of the sensor node

In this section, we review the state of sensor network subsystems, and stipulate the compulsory background in database systems. In subsequent segments, we converse confronts in designing a sensornet database system.

A. Adequate Sensor Network Subsystems in WSN

Examples of sensor devices are starting to arise on the horizon. One class of devices is demonstrated by the mote [23]. Motes contain an 8-bit processor, several megabytes of memory, a low baud-rate radio, and MEMS sensors for detecting temperature, vibration, and ambient light. A class of larger devices [30] comprises spread-spectrum radios, PC-class processors, infrared dipoles, and electret microphones, and acoustic geophones.

Working our way up the layers in Fig. 1, examples of such related research comprise: an well-organized operating system for sensor nodes [23]; low-level network self-configuration structures [17], comprising systems for localizing nodes [31], [32], [33], and performing time organization [11]; a data-centric routing system [25], and possibly cooperative signal administering systems [39], [59] that can, for example, track moving objectives.

B. A Data Models In WSN

In the context of a sensor network, the relational model is best explained as follows. Each sensor creates one or more tuples in network. The node that produces the tuple is called the source.

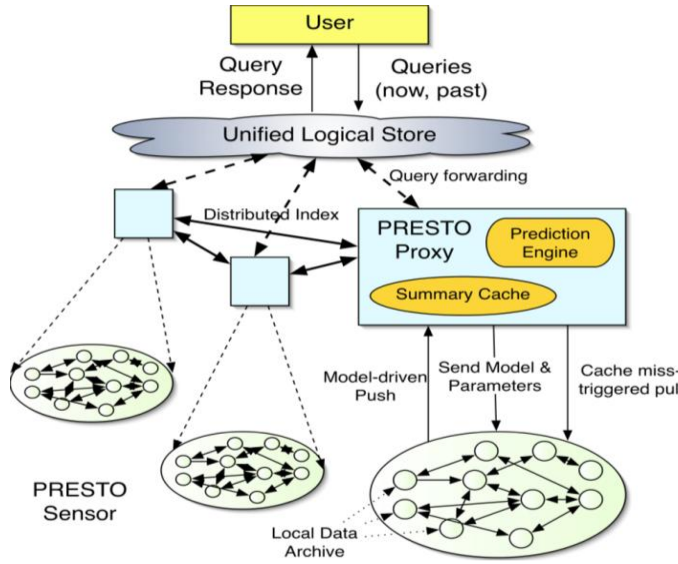

Fig. 3. Sensor Network Software Subsystems

For example, a temperature sensor might generate a tuple of the form <node Location, timestamp, and temperature>. Like wise, at a node that uses aural and vibration signal patterns to sense vehicle, signal processing software might produce a tuple of the form <nodeLocation, timestamp, vehicletype,detectionConfidence> . A compilation of similar tuples from a group of sensors forms a snapshot. In database vocabulary, this snapshot establishes a relational table, which is straight apportioned across the sensors in the unit. Relational tables are predictably accumulated on disks in conservative relational database management systems. Further, it is imperative to note that the tables we deliberate in the sensor network framework are all virtual tables, and they are relational views of the data created by a sensor network. Approaches to these virtual tables are inevitably translated into equivalent data-collecting operations on each appropriate sensor nodes, e.g., Get-Light Intensity, Get Temperature, etc. Virtual tables can be unrestrained, characterizing, for example, streams of data in database.

The objective of the sensornet database proposal should be to ambit location transparency. Specifically outside the database community, the term relational database often suggests notions of strong agreements on storage steadiness and obtainability.

C. Database Operators In WSN Database

The following paragraphs describe some of these traditional database operators in a sensor network context. We stick to an SQL-style multi set semantics; this is typically the desired semantics for aggregation-centric functions [19]. The average temperature on the third floor is an example of an aggregate well defined on a temperature table containing of tuples from sensors in an in-building sensor network system. Most profitable databases afford shared aggregation operators such as SUM, COUNT, AVERAGE, MIN, MAX, and STDDEV.

In conventional databases, the join operator is used to associate data from multiple tables. Further join can be described as a collection over the cross product of a pair of tables; R on S denotes a join of tables R and S. However it is fairly mutual to appliance joins in a more effective fashion that does not for all braces. A common join establish is an equality match across columns of the two tables called as “an equi-join”; e.g., deliberate a temperature table with tuples of the form <nodeLocation, timestamp, and temperature >. It also adopt that some sensor nodes with temperature sensors also have light sensors and each of which creates tuples of the form<nodeLocation, timestamp, lightlevel>. Moreover, an equi-join of these two tables on the nodeLocation column would construct a table with tuples of the form as: <nodeLocation, lightlevel, timestamp, temperature>, where tuples are only described for nodes that have both temperature and light sensors. Fiurther, there are numerous other relational operators like grouping, selection , and projection , difference, union, distinct aggregates ,and duplicate elimination, that we do not converse for conciseness.

III. SENSORNET DATABASE SYSTEM AND FUNDAMENTAL APPROACHES TO DATA MANAGEMENT IN WIRELESS SENSOR NETWORKS

In [24], a scalable and robust communication example, focused dispersion, is suggested. Attribute-value pairs are used to name data created by sensor nodes. A request node sends its attentions of termed data to objective sensor nodes. To progress the performance and save energy, intermediary nodes can cache data and might aggregate the data [27] increases and recovers the focused diffusion method particularly on tests. Tiny OS [16] is a free and open source operating system intended for WSNs. Tiny OS is an entrenched operating system, which is written in the nesC. Hence, Tiny OS can maintenance basic data requirements. Furthermore, users can improve their own applications based on Tiny OS. Tiny DB [28] is a data management system for WSNs constructed on Tiny OS. It can abstract information from WSNs by sending queries. Prominently, Tiny DB permits users designate the data they need to obtain by writing a SQL-like query. Furthermore, for answering a query; Tiny DB desires the data from sensor nodes in the network and directions it backs to a PC (System). In the phase of administering queries filtering and aggregation algorithms might be expended. For example, tree-based routing is used for query distribution, data collection and in-network aggregation the queries. REED [18] spreads Tiny DB with the capacity to procedure joins operations between sensing data and static tables which is built outside the WSN Join operation is accomplished in network approach. Cougar [26] is another dispersed database system to sensor networks that pondered query languages, query optimization, aggregation processing, and multi-query optimization, catalog management [22].

We have said that a sensornet database permits any user to concern a query to the sensor network as if it is a database system and attain a response to that query. Further, there are two apparent recognitions of a sensornet database system. The first one is a centralized i.e., data warehouse recognition, where all data from each node in the network is sent to a selected node within the network involved to which is a large database system. This can be unreasonable in the sensor network perspective since it compels meaningful communication and that needs energy [29]. The other one, a distributed database, can be energy competent when the query rate is less than the rate at which data is produced. This sensornet database architecture respites on two descriptions. The first feature is in-network accomplishment of database operators in a database. When a user postures a query to the network, that query is distributed across the network [56]. Moreover, another work has exposed that in-network administering of sensor data is important to attaining energy-efficient interaction in sensor networks [18].

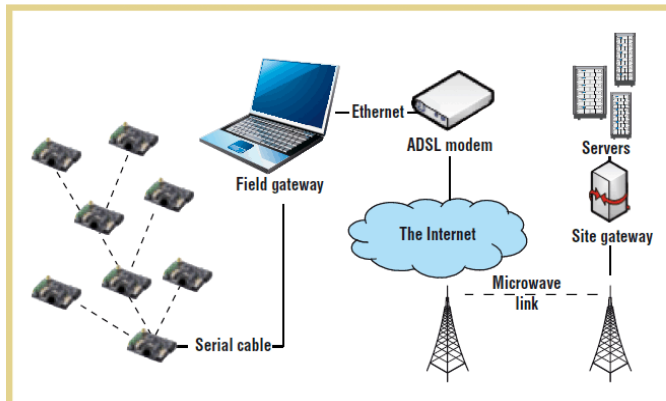

Fig. 4. Sensornet Database

A second feature is that, the sensornet database will deliver estimated outcomes. In sensor networks system, the accessibility of data might be condensed as a consequence of communication loss instigated by impulses in wireless communication or by node failure. This proficiency, called online aggregation, has been suggested in the database literature for large on-line decision support arrangements [22], [20], [43]. The idea of expressing data engendered by sensors, as tuples is casually comparable to the concept of data naming conversed in the context of data-centric sensor network routing [45] and wide-area communication innovation [12]. The COUGAR project at Cornell University [5] is one of the first efforts to model a sensor network as database systems. It incorporates both the SEQ [34] preparation data model and the relational data model by acquainting new operators between sequence data and relational data in a data model. COUGAR does not presently center on abusing the special features of sensor networks nor does it discover the communication between query processing and networking system. Finally, Srivastava et al. [35] Point out the requirement for a data management middleware for sensor network data evaluation and mining, in the context of a particular application. Moreover, this paper takes this a step further and identifies specific challenges in realizing one aspect of this middleware relational database system.

IV. DATA COLLECTION AND OPERATORS USED IN WSN DATA MANAGEMENT

Data collection is extensively used for functions, which gathers all sensed data constantly. In [48], Chu et al. have suggested a mechanism, Ken, using restrictive data transmission to preserve energy by reporting only if the difference between the sensed value and the predicted value is yonder confident limits [7].

A sensor node does not need to report sensed value normally if the predicted error is within the apparently, it is easy to store and access the data at Base station (BS). However, such methods (such as [31], [32], [44]) might be more applicable. As revealed in [33], the prediction models based on the progressive and spatial correlations of data work extremely well in Wireless Sensor Networks. Further, a Model-based withholding is used to deliver uninterrupted data without incessant reporting. In addition, a key problem for, link failure and data suppression is attended. A mobile filtering method for error-bounded data gathering was projected in [26]. Jain et al, have constructed energetic techniques that service maximum filtering of data using a method called stochastic recursive data filtering, to protect possessions subject to conference precision standards [27]. Furthermore, two key issues in Data Management in Sensor Networks are Data Storage and Query Processing. Now we are going to start to highlight some of the research concerns by reflecting the completion of two database operators that are 1). joins and 2). aggregation. Hence, furthermore we'll discuss precisely about these two terminologies.

A. Use of Join in WSN Database Systems

In the sensornet database, the intricacy of the join can differ with the specific query. The easiest example of a join is which joins the temperature and light tables by node location, can be capable nearby. That is, each distinctive node can achieve the join on the temperature and light tuples that it produces before communicating the joined tuple to the query creator [43].

A vibration sensor creates a tuple of the form <eventType, confidenceLevel, vibrationAmplitude,targetLocation>. To compare events from different sensors, one might desire to achieve an equijoin on the eventType column [46].

The database literature has considered numerous generic join implementation methods, such as nested-loop, hash join, and merge-sort [38]. These approaches are blocking. E.g., the hash join algorithms usually used in database systems [18] cannot create any tuples until one of the tables is fully examined. Furthermore, an assortment of non-blocking pipelined join approaches have been established in contemporary years. E.g, is symmetric hash-join [62]. It builds and maintains two hash tables, one for each input table. It is symmetric because the action for each tuple from either table is the identical. A simplification of symmetric hash-joins is the family of join methods called ripple joins [17], and join methods statistically sample the two tables to be combined, in order to construct a stream of joined tuples.

1). Pipelining in a Database

Pipelined joins, because they afford streamed partial answers can empower query improvement. Additionally, pipelining schemes like ripple joins form a low energy methodology to attain estimated answers and can be used collected with sampling [57].

2). Accurate Partitioning and Interactions with Routing

Partitioning and Communications with Routing this methodology points to a procedure used in parallel database systems called partitioning; tuples are partitioned based on their join-column values, and reallocated on the fly across several nodes [10]. The objective here is both to influence parallelism, and to achievement aggregate RAM space across multiple nodes, and the sensor nodes may be memory-constrained. Though, conventional databases did this on a fully connected cluster intersect whereas data-centric storage arrangements are ascendable over random topologies in the wide area. While these approaches are somewhat simplified of relational operators in a sensornet database and can be posed as a routing drawback [53].

B. Use of Aggregation in Database System

The mechanics of computing aggregates is theoretically query is flooded throughout the network or to a quantified geographic section, and the responses are routed on the reverse path trees. Aggregation on multiple nodes is not new and has been broadly discovered in the parallel database information [54].

1). A Taxonomy of aggregates functions

Aggregation on multiple nodes is not new and it was established [16] to classify the different classes of aggregates in terms of their partitioning across numerous nodes in a cluster. In sensor networks, one key enactment objective is to spread the lifetime of the network by reducing communications and, aggregation functions can be helpfully characterized by the sizes of the partial state records that get passed nearby. For example, the AVERAGE aggregate is calculated by each node sending the SUM and COUNT of its readings to its parent, Further with parents sending the SUM of SUM s and COUNT of COUNTs upwards recursively. The root confirms the aggregate by dividing the total SUM by the total COUNT there. Hence the partial state for AVERAGE is two numbers as partial COUNT and partial SUM, and twice the size of the base readings. The first three categorizations were originally offered in the perspective of conventional databases [41]. This is particularly true in communicating settings: user conclusions of information predictors have shown that the first request is often for a big picture of the data that is used to choose what other questions to ask [58].

2). Energy-efficient Aggregation system

The main evidence quality may be appreciably disturbed by packet loss only under convinced circumstances and only for certain kinds of aggregates.

This methodology is appropriate to algebraic aggregates like AVERAGE and has been projected in for online aggregation in conventional databases [49], [50]. In this methodology, tuples in a table are consistently sampled and the resulting average is presumed to signify the actual average. A deviation of this method that is appropriate to counting is a class of probabilistic counting methods that use logarithmic sampling [28, 13]. Recall that the results of a query are sent up the reverse path tree about the originator. For example, a partial SUM to its parent and a node can regularly supply the sum among all nodes within its radio range that are relations of its parent. Hence, we call this method flow-based because it splits up a count or value into many flows and thereby decreases the compassion of the aggregate to damage. Furthermore, to response convinced aggregates like MAX, the query originator could pose a hypothesis answer and see if anybody refutes it; this limits interaction costs to aggregation of refutations [63].

Moreover, we reflect counting based schemes are agreeable to hypothesis testing. Thus, an n -tile is a assumption of the form and there are exactly j nodes j=n readings whose value is greater than a value x R.

Each plan has 3 input tables R, S and T and two join operators conjoining R with S, and S with T. Conclusively, for aggregates where the size of the fractional state is a function of the number of records and data solidity procedures are appropriate.

In the database information, the statistical quality of imprecise outcomes can be robustly designated via confidence intervals for aggregate estimators run over i.i.d Samples of the database to do statistical research on precisely embodying the calculation excellence of consequences.

V. DATA STORAGE AND COMPLEX QUERY OPTIMIZATION PLAN

Several methodologies have been suggested to designate how to store data produced by Wireless Sensor Networks. One category of such storage solutions is that base station collects and stores all data as [17] might be more relevant to answer constant queries. For refining network lifetime, in-network storage procedures have been adopted to resolve ad-hoc queries. These kinds of structures are mainly created on the Data Centric Storage (DCS) concept [33]. In Data Centric Storage, appropriate data are considered and named rendering to its meanings [9]. The major difference among in network Data a Centric Storage scheme is using dissimilar events to sensors recording approaches. The mapping was intended using hash tables in DHT [33] and GHT [34], or expending k-d trees in DIM [35], KDDCS [36,60], and STDCS [11]. STDCS usage sensor location as data indexing instead of the sensed values. Thus, STDCS reports sensors to sensors mapping instead of the readings events to sensors mapping.

STDCS uses a spatiotemporal indexing to equalize query load among sensors. As it's known, indexing methods can suggestively progress the data obtaining query presentation. For Wireless Sensor Networks, another advantage of using index is decreasing cost of data request distribution since the objective of data request can be gotten from A Data Management Tool called ES3N [1] uses Semantic Web procedures to accomplish and query network lifetime for Wireless Sensor Networks with the index. The works in [37] and [35] use a spatially dispersed hashing index procedure to explain range queries in a multidimensional area. The work in [12] suggested a dispersed spatial temporal index structure to trail affecting articles. The work in [38] adopted a time-based index preparation for event query dispensation. This characteristic delivers use accidental to store sketch information to answer queries based on historical data of sensor networks. For boosting query processing presentation and convertible energy, a dispersed index is essential to escort query advancing. In [39], we suggested a in-network antique data storage and query processing scheme based on dispersed indexing [3]. Already It's been designated how database operators might be understood in a sensornet database. For a given query, the order of operator assessment can control resource operation. For energy competence, optimizing complex queries will be a significant objective. As we shall understand, in a sensornet database, a complex query optimization is familiarly connected to routing. To stimulate complex query optimization, reflect a complex join query of the form R on S on T recall that R on S means the join of tables R and S. Joins are commutative and associative, and hence the above expression is correspondent to the expression R on S on T. In the first plan, the join S on T is appraised first and the resulting table is joined with. Further, in the second, the join R on S is assessed first and the resulting table is joined with T. These two query tactics may have dissimilar costs. E.g., if R on S has a small number of tuples, the latter query plan may be more energy effectual than the previous.

The search problem has three parameters: the set of feasible plans, an effectual search algorithm for outcome the minimum cost plan in the space and a cost model for guessing the competence of a plan.

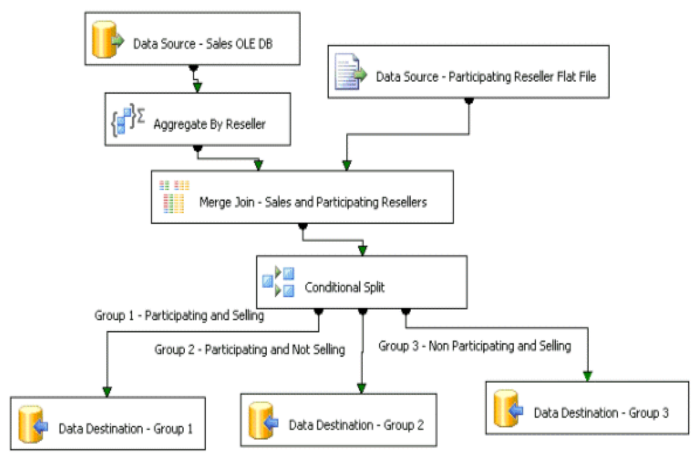

Fig. 5. Complex Query Optimization

Inappropriately, such static plan execution may not be suitable for a sensornet database. Query costs are enormously dynamic in sensor networks [65], [66]. It is also disturbed by network parameters comprising topology, loss rates. Both the data and the interaction in a sensor network are extremely unstable, and hence a more adaptive query optimization methodology is needed [55], [61].

A. *Adaptive Query Optimization Schemes*

In WSNs data gathered from a minidome Sensor Network. The Tiny DB project is based on a query language that supports basic, aggregate, event based, temporal aggregate and even lifetime [2] suggests SPARQL query language For Query Processing, RDF is a focused, marked graph data format for expressive information in the Web system [64].

The results of SPARQL queries can be results sets or RDF graphs. In fact, SPARQL is a Semantic based querying competences. This language backing a range of query types, including monitoring, exploratory, networks health, actuation and offline delivery queries. The consequences of SPARQL queries can be results sets or RDF graphs. In fact, SPARQL is a Semantic Web candidate reference currently SPARQL is entrenched in Jena which is a Java structure for construction Semantic Web applications that offers a programmatic atmosphere for RDF, OWL, and RDFS including a rule-based extrapolation engine [8]. Adaptive query optimization is an area of developing interest in the database community for server side query processing over remote databases [47].

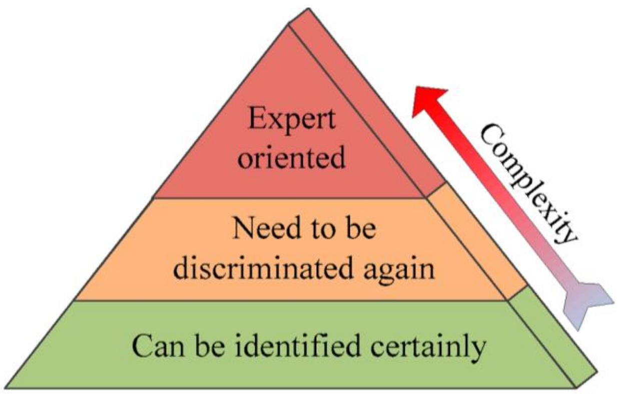

Fig. 6. Adaptive Optimization Schemes

Now briefly we describe about eddie that is mentioned to [4] for more feature. An eddy is a dataflow operator that is interposed between commutative query processing operators. Based on explanations of consumption and manufacture rates of the operators, an eddy routing policy can route received tuples to better operators first, in order to improve the flow of data through all the operators [40], [67]. Thus, eddies energetically do query optimization at runtime. As initially intended for consolidated administering, eddies route data among commutative operators on a single node. Further, another methodology is to have multiple eddies, and mark better global decisions about operator partitioning and placement as labeled in the preceding section [42]. This latter methodology is basically dynamic routing of tuples. The routing protocol is application specific same as the metrics.

This is an attractive example of an incorporation of functionality that would, in more conventional organizations, have been measured as fitting to divisible layers [20].

VI. CONCLUSION

The primary objective of a sensor network is to create worldwide meaningful information from raw local data acquired by specific sensor nodes. Prominently, this purpose must be accomplished in the perspective of persisting as much as conceivable the useful existence of the network and confirming that the network residues highly accessible and endures to deliver precise material in the face of safety occurrences and hardware breakdown. As novel principles based networks are unconfined and low power systems are repeatedly established, we will turn to see the extensive deployment of distributed databases in wireless sensor networks (WSNs). Further, Sensor nodes can be pretend as small computers, tremendously essential in terms of their boundaries and their apparatuses. On these days we can perceive a vide possibility and various applications of distributed database management in wireless sensing procedures. A homogenous query interface for programming data gathering from a wireless sensor network will impressively improve the development has distributed sensing purposes. An important conclusion of research in this area will be an considerate of the proper modularization of sensor network subsystems, and an escalation of the level of combination necessary between distinctive modules to accomplish a strong and competent system, database information has discovered sequential and other sequence centric data models. An example of such a model, which familiarizes sequence-based operators but does not essentially change the implementation and optimization procedures established for the relational model in database system. In a wireless background, data transmission pattern has been predictable, as an operational and ascendable contrivance to distribute frequently demanded evidence to a large number of patrons.

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